

The conic duality theorem

Consider the primal-dual pair

$$\begin{aligned}c^* &= \min c^T x \\ \text{s.t. } &(Ax - b) \in K\end{aligned}$$

PRIMAL

$$\begin{aligned}b^* &= \max p^T b \\ \text{s.t. } &p^T A = c^T \\ &p \in K_*\end{aligned}$$

DUAL

Theorem

- (i) The value $p^T b$ for a **feasible solution** for the dual is **less or equal** to the value $c^T x$ for a **feasible solution** of the primal.
The **duality gap** $c^T x - p^T b \geq 0$ for every **primal-dual pair** (x, p) .
- (ii) If the primal is **strictly feasible** ($Ax - b \succ_K 0$) and bounded below then the dual is feasible and $b^* = c^*$.
- (iii) If the dual is **strictly feasible** and bounded above then the primal is feasible and $b^* = c^*$.
- (iv) Assume one of the problems is bounded and strictly feasible then (x^*, p^*) is an **optimal primal-dual pair**
 - ▶ if and only if $c^T x^* = (p^*)^T b$
 - ▶ if and only if $p^T (Ax - b) = 0$ (**complementarity slackness**)

KKT for convex problems

$$\begin{aligned} & \text{minimize} && f(\mathbf{x}) \\ & \text{s.t.} && \mathbf{g}(\mathbf{x}) \leq \mathbf{0}. \end{aligned}$$

$\hat{\mathbf{x}}$ is a (global) optimal solution if, and only if, there exists a vector $\hat{\mathbf{y}} \in \mathbf{R}^m$ such that

$$(1) \quad \nabla f(\hat{\mathbf{x}}) + \sum_{i=1}^m \hat{y}_i \nabla g_i(\hat{\mathbf{x}}) = \mathbf{0}^T$$

$$(2) \quad \mathbf{g}(\hat{\mathbf{x}}) \leq \mathbf{0},$$

$$(3) \quad \hat{\mathbf{y}} \geq \mathbf{0},$$

$$(4) \quad \hat{\mathbf{y}}^T \mathbf{g}(\hat{\mathbf{x}}) = 0.$$

