

# Introduction to computability

## Tutorial 10

Uncomputability and Complexity

04 December 2018

## Some undecidable languages

*universal language:*  $UL = \{\langle M, w \rangle \mid M \text{ accepts } w\}$

$\overline{UL} = \{\langle M, w \rangle \mid M \text{ rejects or cycles on } w\}$

*halting problem:*  $H = \{\langle M, w \rangle \mid M \text{ stops on } w\}$

*empty-word halting problem:*  $\{M \mid M \text{ stops on } \varepsilon\}$

*existential halting problem:*  $\{M \mid (\exists w) M \text{ stops on } w\}$

*universal halting problem:*  $\{M \mid (\forall w) M \text{ stops on } w\}$

*empty accepted language:*  $\{M \mid L(M) = \emptyset\}$

*recursive accepted language:*  $\{M \mid L(M) \in R\}$

*undecidable accepted language:*  $\{M \mid L(M) \notin R\}$

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3. Let  $M_1$  and  $M_2$  be two Turing machines. Show that the problem that consists of determining if the execution of  $M_1$  on the empty word needs (strictly) fewer steps than the execution of  $M_2$  on the empty word is undecidable.

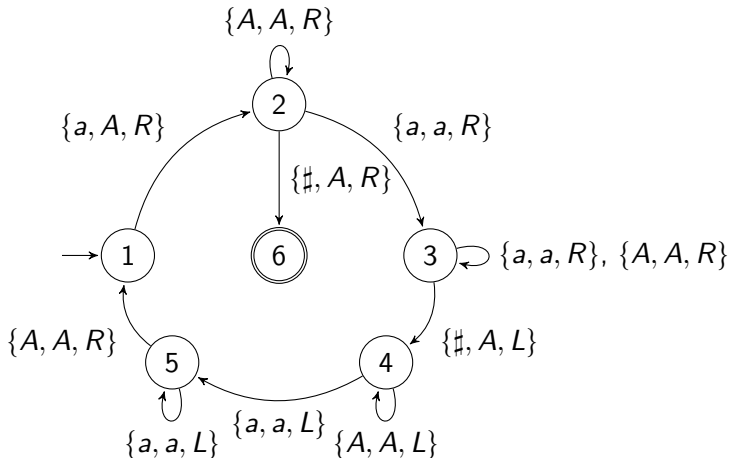
# Complexity

The class **P** is the class of languages *decided* by a *polynomial deterministic* Turing Machine.

The class **NP** is the class of languages that are *accepted* by a *polynomial nondeterministic* Turing Machine.

The class **NPC** is the class of languages that are in NP and that are "hardest to decide".

4. What is the time complexity of the Turing machine below where the entry alphabet is  $\Sigma = \{a\}$ ?



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- ▶ Are the context-free languages in  $P$ ?



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6. Show in a diagram the inclusion relations between  $P$ ,  $NP$ ,  $R$ ,  $RE$ , the regular languages and the context-free languages.

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7. Give a deterministic algorithm for the Hamilton Circuit problem and estimate its complexity.

8. Let  $L_1$  and  $L_2$  be two languages that belong to the same complexity class,  $P$ ,  $NP$  or  $NPC$ . What is the complexity class of the languages  $L_1 \cap L_2$  and  $L_1 \cup L_2$ ?

## Bonus Exercise 11

Give a deterministic algorithm to solve the SAT problem. What is its complexity?