Computation structures

Support for problem-solving lesson #7

Consider the following Java class:

```
public class MyClass {
   public MyClass () {}
   public synchronized void m1() {
      System.out.println("Entering method 1");
      try {Thread.sleep(5000); } catch(Exception e){}
      System.out.println("Exiting method 1");
   }
   public synchronized void m2() {
      System.out.println("Entering method 2");
      try {Thread.sleep(5000); } catch (Exception e){}
      System.out.println("Exiting method 2");
    }
}
```

What could be the outcome of the following programs?

```
MyClass o1, o2;
o1 = new MyClass(); o2 = new MyClass();
new Thread() { public void run() {
    o1.m1(); } }.start();
new Thread() { public void run() {
    o1.m2(); } }.start();
```

```
MyClass o1, o2;
o1 = new MyClass(); o2 = new MyClass();
new Thread() { public void run() {
    o1.m1(); } }.start();
new Thread() { public void run() {
    o2.m1(); } }.start();
```

Both methods m1() and m2() are synchronized.

• Thus, for a given object, several threads cannot access its method simultaneously.

- Consider the first case :
 - One thread wants to execute o1.m1().
 - Another thread wants to execute o1.m2().

First possibility:

- Thread 1 is executed. Entering the function, it aquires the lock on o1.
- From that point, even if Thread 2 wants to execute, it will be blocked outside of m2().
- Thread 1 executes o1.m1() until completion, then releases the lock on o1.
- Thread 2 now can execute o1.m2().

Second possibility:

Same as before, but Thread 2 is executed first, locks Thread 1 outside of m1(), finishes o1.m2(), releases the lock and Thread 1 can execute o1.m1().

• Output:

Entering method 1

<maits 5 seconds>
Exiting method 1
Entering method 2
<maits 5 seconds>
Exiting method 2

or

Entering method 2

<maits 5 seconds>
Exiting method 2
Entering method 1
<maits 5 seconds>
Exiting method 1

- Consider the second case:
 - One thread wants to execute o1.m1().
 - Another thread wants to execute o2.m1().
- The lock is aquired on a given object, thus Thread 1 obtains the lock on o1 while Thread 2 obtains the lock on o2.
- Thread 1 and Thread 2 are thus <u>not</u> in mutual exclusion, even though they both called the same method that was declared <u>synchronized</u>.

• Output:

```
Entering method 1
Entering method 1

<maits ~5 seconds>
Exiting method 1
Exiting method 1
```

An animal shelter has a room to temporarily store animals that transit from their cages to the vet clinic and reversely.

Rules are:

- The room is only used to hold cats or dogs.
- A cat can never enter the room if it already contains a cat or a dog.
- A dog can never enter the room if it already contains a cat.
- There cannot be more than 4 dogs in the room.

Write a solution to this problem using synchronized methods as well as wait(), notify() and notifyAII() calls. Use variables cats and dogs to represent the number of cats and dogs in the room respectively.

Recall:

- wait() will lock the calling thread and release the mutex aquired by synchronized until another thread calls notify() or notifyAll() in a synchronized method of the same object.
- notify() will unlock a single thread that called wait(). The choice of the Thread to unlock is arbitrary and depends on the implementation. In Java 8, the Thread that called wait() first is notified first. If there's none, notify() does nothing. The calling thread can continue its execution.
- notifyAll() will unlock all threads that called wait(), if any. The
 calling thread can continue its execution.

Let's first start without any synchronization.

```
public class Room {
public Room() { }
public
                    void dogEnter()
  dogs++;
                    void dogExit(){
public
  dogs--;
```

```
public
                     void catEnter()
  cats++;
public
                     void catExit(){
  --cats;
int dogs=0; //number of dogs in the room
int cats=0; //number of cats in the room
```

I need to protect cats and dogs, so I use the mutex offered by synchronized.

public class Room { public Room() { } public synchronized void dogEnter() dogs++; public synchronized void dogExit(){ dogs--;

```
public synchronized void catEnter()
  cats++;
public synchronized void catExit(){
  --cats;
int dogs=0; //number of dogs in the room
int cats=0; //number of cats in the room
```

Dogs can only enter if there's less than 4 dogs, and cats can only be alone.

public class Room { public Room() { } public synchronized void dogEnter() throws InterruptedException{ while(dogs>=4 || cats>0) wait(); dogs++; public synchronized void dogExit(){ dogs--;

```
public synchronized void catEnter()
   throws InterruptedException{
  while(cats>0 || dogs>0)
    wait();
 cats++;
public synchronized void catExit(){
  --cats;
}
int dogs=0; //number of dogs in the room
int cats=0; //number of cats in the room
```

When an animal leaves the room, it might let another animal come in, so we must unlock the waiting threads.

public class Room { public Room() { } public synchronized void dogEnter() throws InterruptedException{ while(dogs>=4 || cats>0) wait(); dogs++; public synchronized void dogExit(){ dogs--; notifyAll();

```
public synchronized void catEnter()
   throws InterruptedException{
  while(cats>0 || dogs>0)
    wait();
  cats++;
public synchronized void catExit(){
  --cats;
  notifyAll();
int dogs=0; //number of dogs in the room
int cats=0; //number of cats in the room
```

Why did we use **notifyAll()** instead of **notify()**? Because we could potentially loose a notification.

public class Room { public Room() { } public synchronized void dogEnter() throws InterruptedException{ while(dogs>=4 || cats>0) wait(); dogs++; public synchronized void dogExit(){ dogs--; notifyAll();

```
public synchronized void catEnter()
   throws InterruptedException{
  while(cats>0 || dogs>0)
    wait();
  cats++;
public synchronized void catExit(){
  --cats;
 notifyAll();
int dogs=0; //number of dogs in the room
int cats=0; //number of cats in the room
```

Loosing a notification:

```
Dog 1 wants to enter → OK, dogs = 1, cats = 0
Cat 1 wants to enter → KO, calls wait()
Dog 2 wants to enter → OK, dogs = 2, cats = 0
Dog 3 wants to enter → OK, dogs = 3, cats = 0
Dog 4 wants to enter → OK, dogs = 4, cats = 0
Dog 5 wants to enter → KO, calls wait()
Dog 1 leaves the room → calls notify(), dogs = 3, cats = 0
Cat 1 is awaken by notify(), checks dogs and cats → KO
→ wait()
But Dog 5 could enter and is not notified.
```

Why did we use while instead of if?

Because triggering **notifyAll()** (and even **notify()**) does not guarantee that the condition

that lead to the wait() is now false.

```
public class Room {
public Room() { }
public synchronized void dogEnter()
     throws InterruptedException{
  while(dogs>=4 || cats>0)
    wait();
  dogs++;
public synchronized void dogExit(){
 dogs--;
 notifyAll();
```

```
public synchronized void catEnter()
   throws InterruptedException{
  while(cats>0 || dogs>0)
    wait();
  cats++;
public synchronized void catExit(){
  --cats;
  notifyAll();
int dogs=0; //number of dogs in the room
int cats=0; //number of cats in the room
```

Using if instead of while (even replacing notifyAll() by notify()):

```
    Dog 1 wants to enter → OK, dogs = 1, cats = 0
        Cat 1 wants to enter → KO, calls wait()
        Dog 2 wants to enter → OK, dogs = 2, cats = 0
        Dog 1 leaves the room → calls notify(), dogs = 1, cats = 0
        Cat 1 is awaken by notify(), and resumes its execution → dogs = 1, cats = 1 → KO
```

A bank asks your help to develop a Java program that performs the payments.

Bank accounts are stored in objects of class **Account** that advertise three non-atomic methods:

- void credit(double amount) to credit amount to the account,
- void debit(double amount) to debit amount from the account, and
- String getIBAN() to get the IBAN of the account.

You must write a method called

transfer(Account from, Account to, double amount)

that will be used in the context of multi-threading, and ensure synchronization is performed in such a way as to keep the accounts in a coherent state while avoiding deadlocks.

Recall

You can use synchronized in three ways:

- For a whole method
 - e.g. synchronized int theMethod() { /* mutex on the method*/}
- For a block inside a method (on the current object)
 - e.g. synchronized(this) { /* mutex for this block only */}
- For a block inside a method (on a different object)
 - e.g. synchonized(obj) { /*mutex aquired for object obj*/ }

First try

```
public synchronized void transfer(Account from, Account to, double amount)
{
    from.debit(amount);
    to.credit(amount);
}

Is it good? If not, why?
The mutex is aquired on the whole object.
```

I cannot work on account3 and account4 if I'm already working on account1 and account2.

→ Too restrictive

```
Second try
 public void transfer(Account from, Account to, double amount)
    synchronized(from) {
        synchronized(to) {
           from.debit(amount);
           to.credit(amount);
Is it good? If not, why?
Imagine the following:
Two transfer orders arrive, one from account1 to account2, and one from
  account2 to account1
Thread1 gets the lock on account1, then Thread2 gets the lock on account2,
```

hence Thread1 cannot get the lock on account2 and Thread2 cannot get the

→ deadlock!

lock on account 1.

Third try

To solve, the deadlock problem, we must find an order for the resources, and lock them accordingly;

```
if(from.getIBAN().compareTo(to.getIBAN()) < 0) {
    synchronized(from) {
        synchronized(to) {
            from.debit(amount);
            }
        } else {
        synchronized(to) {
            synchronized(from) {
                from.debit(amount);
            }
        }
}</pre>
```

public void transfer(Account from, Account to, double amount)

Transfer from BE43 0000 0000 0101 to BE32 0000 0000 0202 (Thread1)

Transfer from BE32 0000 0000 0202 to BE43 0000 0000 0101 (Thread2)

BE32... < BE43... → lock on BE32... for both threads

→ No deadlock